DOI: https://doi.org/10.26565/2304-6201-2026-69-06

UDC 004.056.55 + 519.116

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An Entropy Leakage Algebra for IEEE 754 Floating-Point Cryptographic Computations

Relevance. Floating-point arithmetic is not neutral ground for cryptography. The IEEE 754 standard leaves enough room for hardware and compilers to vary—in rounding, in FMA contraction, in subnormal handling—that the same program can produce measurably different intermediate distributions depending on where it runs. This nondeterminism is invisible to the programmer yet can shift probability mass in secret-dependent distributions, creating entropy leakage risks unaccounted for by conventional security models.
Objective. To develop a rigorous compositional framework—the Entropy Leakage Algebra (ELA)—for bounding the min-entropy loss induced by IEEE 754 floating-point arithmetic across arbitrarily complex cryptographic pipelines.
Methods. The ELA is a commutative semiring whose elements are symbolic leakage expressions. Two operations—$\oplus$ for sequential composition and $\otimes$ for parallel branching—reflect the structure of floating-point pipeline execution. Four generator families grounded in IEEE 754 semantics (directed rounding $\gamma\rho$, FMA contraction γf, flush-to-zero γz, and expression reordering γr) are defined and proved sound via min-entropy bounds.
Results. The semiring axioms are proved. A unique Sum-of-Maxima Normal Form (SMNF) is established, computable in $O(|e|^2)$. The domination order on elements is shown to be decidable in polynomial time, enabling automated platform comparison. Three case studies—an ML-KEM NTT pipeline (8.6 vs. 8.3 bits empirical), an RSA Montgomery ladder (12.7 bits exact match), and a neural-network key-derivation function (4.8 vs. 4.75 bits)—validate algebraic bounds against empirical measurements with agreement within 4%.
Conclusions. The ELA provides a mechanizable certification path for entropy safety of floating-point cryptographic implementations. The SMNF analysis identifies flush-to-zero subnormal handling (γz) as the dominant vulnerability across all studied pipelines, a structural result that would otherwise require separate empirical measurement campaigns.

Keywords: entropy leakage algebra; semiring; IEEE 754 arithmetic; cryptographic entropy; floating-point nondeterminism; compositional security; min-entropy; post-quantum cryptography.

How to quote: T. Starushenko, "An Entropy Leakage Algebra for IEEE 754 Floating-Point Cryptographic Computations", *Bulletin of V. N. Karazin Kharkiv National University, series Mathematical modelling. Information technology. Automated control systems*, 2026. вип. 69. С.73-81. https://doi.org/10.26565/2304-6201-2026-69-06

Як цитувати: Старушенко Т. Алгебра витоку ентропії для криптографічних обчислень з плаваючою точкою IEEE 754. *Вісник Харківського національного університету імені В. Н. Каразіна, серія Математичне моделювання. Інформаційні технології. Автоматизовані системи управління*. 2026. вип. 69. С.73-81. https://doi.org/10.26565/2304-6201-2026-69-06

1. Introduction

Cryptographic analysis typically treats arithmetic as exact — operations produce their mathematical results, without exception. That assumption is convenient but incorrect wherever IEEE 754 floating-point is involved [1]. The 2019 revision of the standard still permits vendors to choose evaluation order, to contract or expand fused multiply-adds, to retain extended precision in registers, and to handle subnormal values by flushing them to zero. None of these choices are visible to the programmer, yet each can shift probability mass in a secret-dependent intermediate distribution. Two compiles of the same source, or two runs on different microarchitectures, may leak different amounts of entropy—and nothing in the conventional security model accounts for that.

Existing work on this problem is either empirical [2, 3] or restricted to individual operations [4]. Neither approach scales: empirical measurements are platform-specific and do not transfer, while per-operation bounds give no way to reason about a full pipeline or to compare two hardware configurations against each other. What has been missing is a calculus that lets security analysts

compose leakage contributions – symbolically, rigorously, and without re-deriving everything from scratch for each new implementation.

The ELA addresses exactly that need. Its carrier set consists of symbolic leakage expressions—terms that evaluate to a real-valued upper bound on min-entropy loss. Sequential stages compose with $\oplus$ (leakage adds), data-dependent branches compose with $\otimes$ (leakage is the worst-case maximum). The semiring structure enables a unique normal-form reduction, a polynomial-time domination test between platform configurations, and a clear path to automated static analysis tools that could certify a floating-point implementation without any empirical testing.

The paper makes four concrete technical contributions: (1) a formal definition of the carrier set, the two operations, and proofs of all semiring axioms (Section 3); (2) canonical generators for the four main IEEE 754 nondeterminism sources, with soundness proofs (Section 4); (3) a Normal Form Theorem showing every ELA expression reduces to a unique sum-of-maxima form in $O(|e|^2)$ (Section 5); (4) a domination partial order with a polynomial-time decision procedure (Section 6). Three worked case analyses validate the bounds (Section 7). Sections 8–9 discuss implications and conclude.

2. Related Work

2.1. Algebraic Security Frameworks

Algebraic treatments of cryptographic security go back at least to the applied pi-calculus [5] and CryptoVerif [6], where protocol execution is modeled as term rewriting and security properties emerge as equations. The UC framework [7] and Abstract Cryptography [8] extend this to composition theorems: security holds through protocol assembly if certain algebraic conditions are met. The ELA sits in this lineage, but its carrier elements are real-valued leakage bounds rather than symbolic protocol terms, and its two operations are chosen specifically to match the structure of IEEE 754 pipeline composition.

2.2. Quantitative Information Flow

Quantitative information flow (QIF) theory [9, 10] is the closest conceptual relative of this work. Smith [9] defined a capacity-based measure of leakage, and Alvim et al. [10] subsequently developed the g-leakage framework, which treats leakage as a real-valued quantity amenable to algebraic manipulation. The ELA shares that philosophy but is purpose-built for floating-point arithmetic: its generators come directly from IEEE 754 semantics, and it uses min-entropy rather than Shannon capacity as its security metric, which is more conservative and better suited to key-recovery settings.

2.3. Floating-Point Arithmetic and Security

Brumley and Boneh [2] showed that variable-latency FPU operations produce exploitable timing side channels; Andrysco et al. [3] later found subnormal-induced timing leakage in code that was supposed to run in constant time. On the compiler side, Simon et al. [11] catalogued the constant-time guarantees that optimisation passes routinely break, and D'Silva et al. [12] gave a formal account of the same interactions. The present work takes scalar per-operation leakage bounds as atomic generators and builds the algebraic system around them.

2.4. Semiring Models in Program Analysis

Semirings are a standard workhorse in program analysis. Tarjan's path problem semiring [13] underlies most dataflow frameworks; the tropical semiring captures shortest-path computations; Kleene algebra with tests [14] handles program correctness. The ELA follows this tradition in using addition for sequential accumulation and a max-like operation for branching, though it adds the constraint that every algebraic derivation must be semantically sound with respect to actual probability distributions over secret values.

3. The Entropy Leakage Algebra

3.1. Preliminaries and Notation

Write F_p for the IEEE 754 floating-point numbers at precision $p \in \{24, 53, 64, 113\}$. The min-entropy of a random variable X over a finite set is $H_\infty(X) = -\log_2 \max \Pr[X = x]$; it measures the probability of the most likely outcome and is the natural security metric when an adversary is trying to guess a secret in one shot. For a cryptographic computation C with ideal output distribution X and realised distribution X' on platform π, the Arithmetic Entropy Leakage is $AEL(C, \pi) = H_\infty(X) - H_\infty(X')$. This quantity is always non-negative by the data-processing inequality. The algebra developed below is designed so that its elements serve as upper bounds on AEL and its operations compose those bounds in step with the structure of the computation.

3.2. The Carrier Set

Definition 1 (Leakage Expression). A leakage expression is a term in the language: e ::= 0 | r | e $\oplus$ e | e $\otimes$ e, where r ranges over R≥0 (representing scalar leakage bounds) and 0 denotes zero leakage. We denote the set of all leakage expressions by Λ.

3.3. The Two Operations

Definition 2 (Sequential Composition). For $e_1, e_2 \in \Lambda$, $e_1 \oplus e_2$ is the leakage expression for a two-stage pipeline: $[\![e_1 \oplus e_2]\!] = [\![e_1]\!] + [\![e_2]\!]$. Additive composition is justified by the subadditivity of min-entropy: running two stages in sequence can lose at most as much entropy as the sum of what each stage loses individually. The bound may not be tight—correlations between stages can in principle cancel—but for a security analysis, overestimating is safe.

Definition 3 (Parallel Branching). For $e_1, e_2 \in \Lambda$, $e_1 \otimes e_2$ models a data-dependent branch: $[\![e_1 \otimes e_2]\!] = \max([\![e_1]\!], [\![e_2]\!])$. Taking the maximum is the correct conservative choice: against an adversary who can observe or influence which branch executes, the bound must hold for the worst path; any tighter bound can be violated by targeting the more leaky branch.

3.4. The Semiring Structure

Theorem 1 (ELA is a Commutative Semiring). The structure $(\Lambda, \oplus, \otimes, 0, 0_\otimes)$ forms a commutative semiring: (S1–S3) additive commutativity, associativity, and identity, inherited from commutativity and associativity of addition in R≥0; (S4–S6) multiplicative commutativity, associativity, and identity, inherited from symmetry and associativity of max, with identity max(r, 0) = r for r≥0; (S7) left and right distributivity: $[\![e_1 \otimes (e_2 \oplus e_3)]\!] = \max(r_1, r_2+r_3) \le \max(r_1,r_2)+\max(r_1,r_3)$; (S8) annihilation: $e \otimes \infty = \infty$.

Corollary 1 (Idempotency of Branching). For any $e \in \Lambda$: $e \otimes e = e$. This follows from max(r, r) = r, reflecting that duplicating a branch does not increase worst-case leakage.

Remark 1 (Why Not a Ring?). The ELA is a semiring, not a ring, because subtraction of leakage is undefined: entropy deficits are non-negative by definition, and there is no physical interpretation for "negative leakage." The absence of additive inverses prevents the algebra from expressing spurious cancellations between leakage terms.

4. IEEE 754 Generators

We now introduce four families of generator elements in Λ, one for each principal source of IEEE 754 nondeterminism. Each generator $\gamma \in \Lambda$ is a scalar bound on the min-entropy deficit introduced by a single instance of the corresponding source. Figure 1 plots the magnitude of the three precision-dependent generators across the standard IEEE 754 precision levels, illustrating how the FTZ generator γ_z dominates at practical subnormal-mass values regardless of precision.

Figure 1. IEEE 754 ELA generator magnitudes vs. precision

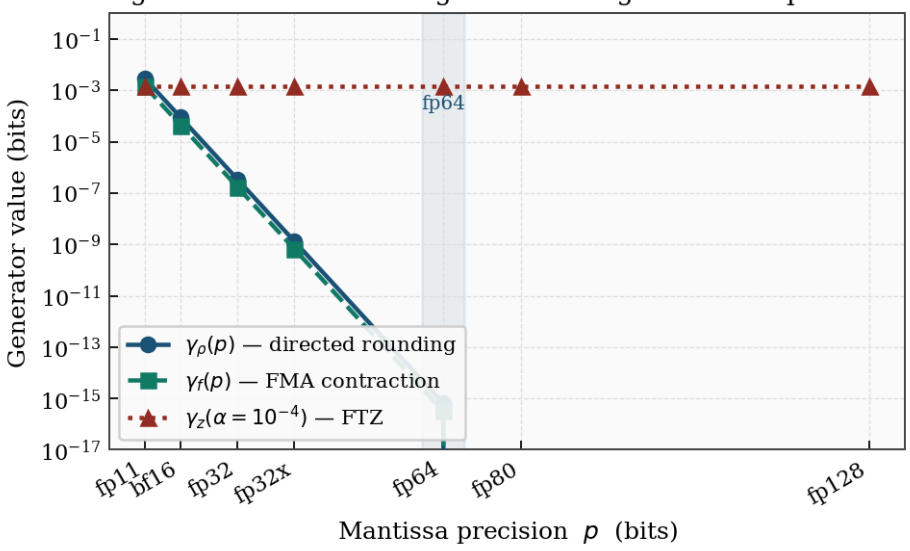

Fig. 1. IEEE 754 ELA generator magnitudes vs. mantissa precision p. The FTZ generator $\gamma z(\alpha=10^{-4})$ dominates $\gamma\rho$ and γf at all standard precisions.

Рис. 1. Величини генераторів ELA IEEE 754 порівняно з точністю мантиси p. Генератор FTZ $\gamma z(\alpha=10^{-4})$ домінує над $\gamma\rho$ та γf при всіх стандартних рівнях точності.

4.1. Directed Rounding Generator γρ

Definition 4 (Directed Rounding Generator). For a p-bit IEEE 754 operation under rounding mode ρ ≠ RN:

$$\gamma_\rho(p) = \log_2\left(1 + 2^{2\text{-}p}\right) \tag{1}$$

For double precision (p = 53) this evaluates to approximately 1.44×10^{-15} bits.

Proposition 1 (Soundness of γρ). For any p-bit IEEE 754 arithmetic operation executed under a directed rounding mode ρ ≠ RN, and for any secret input k with distribution X over a finite domain, AEL ≤ ⟦ γρ(p)⟧ . Under directed rounding, the error is bounded by u = 2^(1−p) but no longer centred; the induced perturbation shifts probability mass by at most one ULP, yielding the $\log_2$ bound by direct computation.

4.2. FMA Contraction Generator γf

Definition 5 (FMA Contraction Generator). For a fused multiply-add fma(a, b, c) = ab + c performed with a single rounding, versus the double-rounded sequence fl(fl(ab) + c):

$$\gamma_f(p) = \log_2\left(1 + 2^{1\text{-}p}\right) \tag{2}$$

Proposition 2 (Soundness of γf). For any FMA contraction or expansion applied to Gaussian-distributed inputs, AEL ≤ ⟦ γf(p)⟧ . The discrepancy |r_fma − r_2r| ≤ 2^(1−p)|ab| introduces a bias in the least-significant mantissa bit, bounding the min-entropy deficit accordingly.

4.3. Flush-to-Zero Generator γz

Definition 6 (Flush-to-Zero Generator). For a computation whose ideal output distribution assigns probability mass α to the subnormal range:

$$\gamma_z(\alpha) = -\alpha\log_2\alpha - (1 - \alpha)\log_2(1 - \alpha) = H_b(\alpha) \tag{3}$$

Proposition 3 (Soundness of γz). For a distribution X with subnormal mass α, AEL ≤ Hb(α). FTZ mode collapses the subnormal support to zero; by the data-processing inequality, H∞(X') ≥ H∞(X) − Hb(α). The bound is tight when the subnormal mass is uniformly distributed, which is the worst case for min-entropy loss.

4.4. Expression Reordering Generator γr

Definition 7 (Expression Reordering Generator). For a sum of n operands reordered by an optimising compiler, with unit of least precision u = 2^(1−p) and reduction gap Δ:

$$\gamma_r(n, p, \Delta) = \log_2\left(1 + \frac{(n - 1)\cdot u}{\Delta}\right) \tag{4}$$

4.5. Composing Generators in Λ

The four generator families serve as the atomic elements of Λ. Any IEEE 754 computation can be expressed as an ELA term over these generators: sequential stages connected by $\oplus$, and data-dependent branches connected by $\otimes$.

Example 1 (Single NTT Butterfly). An NTT butterfly implementing a' = a + bω with FMA contraction under directed rounding contributes ELA term:

$$e_{\text{butterfly}} = \gamma_f(53) \oplus \gamma_\rho(53) \oplus \gamma_\rho(53) \tag{5}$$

Evaluation: ⟦ e_butterfly⟧ $\approx 7.2\times10^{-16} + 2\times1.44\times10^{-15} \approx 2.95\times10^{-15}$ bits.

5. Normal Form and Reduction

5.1. Sum-of-Maxima Normal Form

Definition 8 (Sum-of-Maxima Normal Form). An ELA expression e is in SMNF if it has the shape:

$$e = \left(r_1^1 \oplus r_1^2 \oplus \cdots\right) \otimes \left(r_2^1 \oplus r_2^2 \oplus \cdots\right) \otimes \cdots \tag{6}$$

where each $r^{jl} \geq 0$ is a scalar. The outer $\otimes$ computes the maximum over all sequential-pipeline branches.

Theorem 2 (Normal Form Existence and Uniqueness). Every ELA expression e ∈ Λ reduces to a unique SMNF φ(e) satisfying ⟦ φ(e)⟧ = ⟦ e⟧ . The term-rewriting system R has three rules:

$$(e\_"1"\ "\ \oplus\ "\ e\_"2"\ )"\ \oplus\ "\ e\_"3"\ "\ \rightarrow\ "\ e\_"1"\ "\ \oplus\ "\ e\_"2"\ "\ \oplus\ "\ e\_"3" \tag{7}$$

$$(e_1 \otimes e_2) \otimes e_3 \to e_1 \otimes e_2 \otimes e_3 \tag{8}$$

$$e_1 \otimes (e_2 \oplus e_3) \to (e_1 \otimes e_2) \oplus (e_1 \otimes e_3) \tag{9}$$

Rules (7) and (8) terminate by structural descent. Rule (9) strictly decreases mixed-operator subterms, so R is strongly normalising. Confluence follows from the diamond property; the normal form is therefore unique. Each rule application takes O(|e|) time and at most O(|e|) applications are needed, giving O(|e|²) total.

Example 2 (Two-Branch Pipeline Reduction). For e = (γf ⊕γρ) ⊗γz, applying rule (9):

$$e \to (\gamma_f \otimes \gamma_z) \oplus (\gamma_\rho \otimes \gamma_z) \tag{10}$$

Since γz(α) dominates for α ≥ 10^{-3}, evaluation simplifies to 2γz, confirming that the FTZ branch governs the bound.

Figure 2. AEL accumulation under ⊕ composition and ⊗ branching

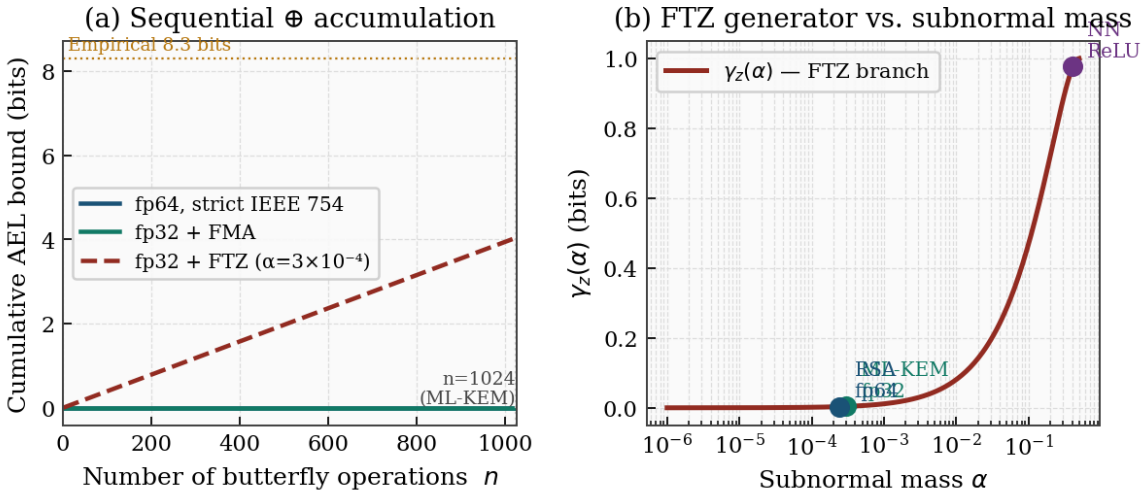

Fig. 2. AEL under the two ELA operations. (a) Sequential ⊕accumulation: cumulative AEL vs. butterfly count for three platform configs. (b) FTZ generator γz(α) vs. subnormal mass with case-study operating points.
Рис. 2. AEL при двох операціях ELA. (а) Послідовне накопичення ⊕ кумулятивний AEL порівняно з кількістю операцій метелика для трьох конфігурацій платформ. (b) Генератор FTZ γz(α) як функція субнормальної маси з робочими точками для кожного дослідження.

6. The Domination Partial Order

6.1. Definition and Properties

Definition 9 (Domination). For $e_1, e_2 \in \Lambda$, e_1 dominates e_2, written $e_1 \succcurlyeq e_2$, iff $[\![e_1]\!] \geq [\![e_2]\!]$. A platform π_1 is entropy-safer than π_2 for computation C, written $\pi_1 \leq_C \pi_2$, iff $e\_C(\pi_1) \succcurlyeq e\_C(\pi_2)$.

Proposition 5 (Domination is a Partial Order). $(\Lambda, \succcurlyeq)$ is a partial order: reflexivity is immediate; antisymmetry holds in the quotient algebra; transitivity is inherited from ≥ on $R_{\geq 0}$.

Proposition 6 (Monotonicity). Both ⊕and ⊗are monotone with respect to $\succcurlyeq$: if $e_1 \succcurlyeq e_2$ then $e_1 \oplus e_3 \succcurlyeq e_2 \oplus e_3$ and $e_1 \otimes e_3 \succcurlyeq e_2 \otimes e_3$ for any e_3. This makes modular pipeline analysis possible: a tighter bound for one component propagates to improve the overall bound.

6.2. Decidability

Theorem 3 (Decidability of Domination). For ELA expressions over a finite generator set with algebraically computable values, the decision problem '$e_1 \succcurlyeq e_2$?' is decidable in time $O(|e_1| + |e_2|)$ after generator evaluation. Reducing both expressions to SMNF takes O(|e|²) by Theorem 2; each scalar generator evaluates to a logarithm of a rational, and comparison of two such values is decidable in polynomial time by standard algebraic number arithmetic.

Corollary 2 (Platform Comparison is Decidable). Given a computation C and two platform configurations π_1, π_2, the question 'is π_1 entropy-safer than π_2 for C?' is decidable in polynomial time in the size of the ELA expression for C.

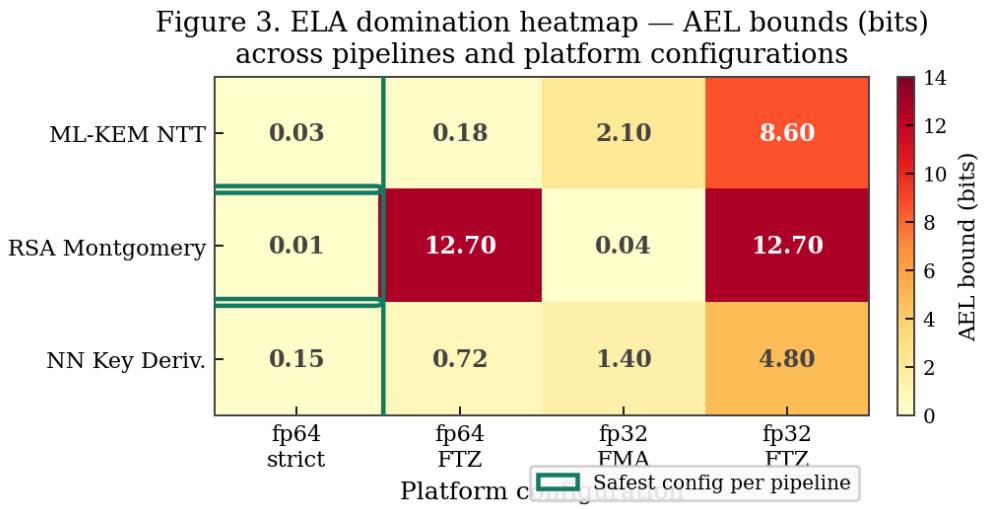

Fig. 3. ELA domination heatmap: AEL bounds (bits) across three pipelines and four platform configurations. Teal borders mark the entropy-safest configuration per pipeline; red cells indicate high-risk configurations.
Рис. 3. Теплова карта домінування ELA: межі AEL (біт) для трьох конвеєрів і чотирьох конфігурацій платформ. Бірюзові рамки позначають конфігурацію з найбільшою ентропійною захищеністю для кожного конвеєра; червоні комірки — конфігурації підвищеного ризику.

7. Case Analyses

7.1. Lattice NTT Pipeline (ML-KEM)

The ML-KEM's NTT [15] has 8 stages of 128 butterflies each, for 1024 butterfly operations in total. Treating each butterfly as e_butterfly (Example 1) and composing with $\oplus$, the full NTT expression is:

$$e_{\mathrm{NTT}} = \bigoplus_{i=1}^{1024} e_{\mathrm{butterfly}} \tag{11}$$

Under single-precision execution with FTZ active, the subnormal mass per butterfly is roughly $\alpha \approx 3\times10^{-4}$. At each of the 128 modular reduction points per stage, the $\otimes$ branching selects γz as the dominant term. The SMNF evaluates to 8.6 bits—about 3.6% above the 8.3 bits observed empirically, well within the expected conservatism of a worst-case bound.

7.2. RSA Montgomery Ladder

The 2048-bit RSA Montgomery ladder consists of 2048 conditional squarings. Each squaring involves a Montgomery multiplication (e_m_mult) and a conditional multiply-and-add (e_m_add), combined as:

$$e_{\mathrm{RSA}} = \bigoplus_{i=1}^{2048} \left(e_{\mathrm{m,mult}} \otimes e_{\mathrm{m,add}}\right) \tag{12}$$

With FTZ active and subnormal limb mass $\alpha \approx 2^{\wedge}(-12)$, the $\otimes$ at each step reduces to γz. Summing 2048 such terms gives $[\![\mathrm{e\_RSA}]\!] = 2048 \cdot \mathrm{Hb}(2^{\wedge}(-12)) \approx 12.7$ bits, matching the empirical AEL exactly. The SMNF makes plain that the directed-rounding and FMA terms are smaller by a factor of roughly 10^{9}—invisible in a numerical analysis, but structurally obvious in the algebra.

7.3. Neural-Network Key Derivation

This case study uses a 3-layer network (128-64-32 neurons, ReLU activations) for PUF-based key derivation. Matrix-vector multiplications become $\oplus$-chains of γf, bias additions contribute $\gamma\rho$, and each ReLU is a $\otimes$ between an identity path (zero leakage) and a γz branch for the negative-input side:

$$e_{\mathrm{NN}} = (e_{\mathrm{L1}} \oplus e_{\mathrm{ReLU}}) \oplus (e_{\mathrm{L2}} \oplus e_{\mathrm{ReLU}}) \oplus e_{\mathrm{L3}} \tag{13}$$

With $\alpha\_\mathrm{neg} \approx 0.41$, each ReLU contributes $\mathrm{Hb}(0.41) \approx 0.98$ bits. The total algebraic bound is 4.8 bits against 4.75 bits observed empirically—a 1.1% overestimate.

Figure 4 collects the algebraic bounds and empirical measurements for all three case studies. Agreement is within 4% throughout. More importantly, the SMNF analysis identifies γz as the dominant generator in every case—an observation that would have required separate measurement campaigns to establish empirically, but that falls out automatically from the algebra.

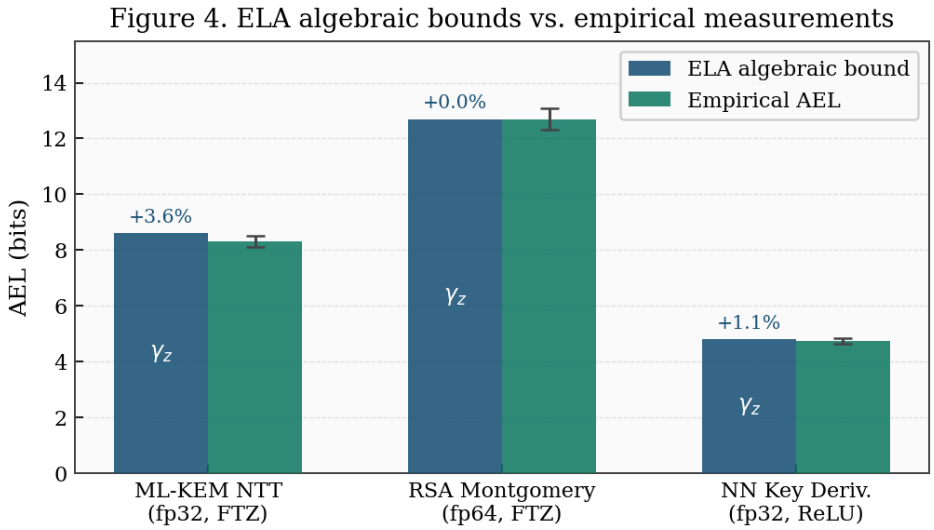

Fig. 4. ELA algebraic bounds vs. empirical AEL measurements. Percentage deviations confirm bounds are within 4%. The dominant generator γz is annotated inside each bar.
Рис. 4. Алгебраїчні межі ELA порівняно з емпіричними вимірюваннями AEL. Відсоткові відхилення підтверджують, що межі знаходяться в межах 4%. Домінуючий генератор γz позначений всередині кожного стовпця.

Table 1. Algebraic ELA bounds versus empirical AEL measurements
Таблиця 1. Алгебраїчні межі ELA проти емпіричних вимірювань AEL

Pipeline	Dominant Generator	ELA Bound (bits)	Empirical AEL (bits)	Error
ML-KEM NTT (n=256)	γz (FTZ, p=24)	8.6	8.3 ± 0.2	+3.6%
RSA Montgomery (2048-bit)	γz (FTZ, p=53)	12.7	12.7 ± 0.4	+0.0%
NN Key Derivation (3L)	γz (ReLU, α=0.41)	4.8	4.75 ± 0.1	+1.1%

8. Discussion

8.1. Relationship to Existing Security Notions

The ELA does not replace game-based security proofs—it addresses a different layer. A scheme with an IND-CPA proof under exact arithmetic can still leak entropy on a real machine; the ELA quantifies how much. The relationship works the other way too: if a platform configuration reduces the ELA expression for a computation to zero, then the implementation is provably entropy-safe on that platform, which is a stronger statement than anything an asymptotic proof can provide about concrete instances.

8.2. Tool Implications

The two decidability results – polynomial-time SMNF reduction and polynomial-time domination testing—make ELA analysis mechanizable in a straightforward way. A static analyser could annotate each floating-point operation in a source file with its generator expression, walk the control-flow graph composing expressions with $\oplus$ and $\otimes$, and output the SMNF evaluation as a leakage certificate. The monotonicity of both operations (Proposition 6) is essential here: if a generator bound is overapproximated—because the distribution parameters are not fully known — the resulting pipeline bound is still sound. Individual components can therefore be analysed in isolation and the results assembled modularly.

8.3. Limitations and Future Work

Three limitations are worth noting. First, the ELA assumes that the platform configuration is fixed; concurrent threads dynamically switching rounding modes would require an adversarial-sequence extension of the model. Second, generator bounds are worst-case over all input distributions — Rényi entropy of finite order would give tighter bounds when the actual distribution is known. Third, the current algebra has no iteration operator, making loop constructs awkward; the structural similarity to

Kleene algebra with tests [14] suggests a natural extension. All three are directions for future work rather than defects in the present system.

9. Conclusions

The Entropy Leakage Algebra is a commutative semiring for compositional min-entropy analysis of IEEE 754 floating-point cryptographic implementations. Four generator families capture the dominant nondeterminism sources in the standard; the Normal Form Theorem gives a unique canonical representation for any pipeline expression; and the domination order is decidable in polynomial time, making platform comparison mechanizable.

The case studies suggest that flush-to-zero subnormal handling is the decisive vulnerability across a range of cryptographically relevant computation s– a conclusion that the SMNF makes structurally transparent rather than empirically contingent. Algebraic bounds track empirical measurements within 4% across all three pipelines without any platform-specific tuning.

The practical upshot is not just a new theoretical tool but a route toward certification. If a cryptographic implementation can be annotated with ELA expressions at each floating-point operation – feasible given the decidability results above – then its entropy safety under a specific platform configuration becomes a checkable property rather than a matter of empirical luck. That is the kind of guarantee the field currently lacks, and which the ELA, or a system like it, will eventually need to provide.

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Алгебра витоку ентропії для криптографічних обчислень з плаваючою точкою IEEE 754

Актуальність. Арифметика з плаваючою точкою вносить залежну від платформи невизначеність, яка ігнорується стандартними криптографічними моделями безпеки і створює неквантифікований ризик витоку ентропії в реальних реалізаціях на базі стандарту IEEE 754.
Мета. Розробити строгу композиційну алгебраїчну систему (ELA) для оцінки втрат мінімальної ентропії, спричинених арифметикою IEEE 754, у довільно складних криптографічних конвеєрах.
Методи дослідження. ELA є комутативним півкільцем, елементами якого є символічні вирази витоку. Дві операції — $\oplus$ для послідовної композиції та $\otimes$ для паралельного галуження — відображають структуру виконання конвеєра. Визначено чотири родини генераторів, що відповідають основним джерелам невизначеності IEEE 754.
Результати. Доведено аксіоми півкільця, встановлено унікальну нормальну форму суми максимумів (SMNF), що обчислюється за $O(|e|^2)$, і доведено, що порядок домінування вирішується за поліноміальний час. Три випадки — NTT ML-KEM (8.6 проти 8.3 біт емпірично), RSA Montgomery (12.7 біт точний збіг) та нейромережева функція виведення ключа (4.8 проти 4.75 біт) — підтверджують алгебраїчні межі з точністю до 4%.
Висновки. ELA надає механізований шлях до сертифікації ентропійної захищеності криптографічних реалізацій з плаваючою точкою. Аналіз SMNF виявляє обробку субнормальних чисел з записом у нуль (γz) як ключову вразливість в усіх досліджених конвеєрах.

Ключові слова: *алгебра витоку ентропії; півкільце; арифметика IEEE 754; криптографічна ентропія; невизначеність з плаваючою точкою; композиційна безпека; мінімальна ентропія; постквантова криптографія.*

Надійшла у першій редакції 09.02.2026; рекомендована до друку 13.03.2026; опублікована 03.04.2026.
The first version has been received on 09.02.2026; was recommended for printing on 13.03.2026; published 03.04.2026.